Provably Secure Scalable Distributed Authentication for Clouds

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- Cloud computing is a significant technology
- There are some challenges and problems in cloud providers::
 - Servers are compromised and the secrets are vulnerable to theft.
 - Secure user authentication is an important issue of cloud services

If authentication is breached, confidentiality and integrity of the data or services may be compromised

OpenStack Identity service (for clouds):

- username and password
- Lightweight Directory Access Protocol
- Kerberos
- TLS/SSL

Incidents and issues

- Weak passwords (dictionary and rainbow table attacks)
- Poor maintenance and implementation (side channel attacks)
- OneLogin
- Golden Ticket Attack

Multi-server environment

• Brainard, J., Juels, A., Kaliski, B., Szydlo, M.: A new two-server approach for authentication with short secrets.

Password-authenticated key exchange

- Boyen, X.: HPAKE: password authentication secure against cross-site user impersonation.
- Katz, J., Ostrovsky, R., Yung, M.: Efficient password-authenticated key exchange using human-memorable passwords.

Treshold solutions

- Isler, D., Küpcü, A: Distributed Single Password Protocol Framework.
- Di Raimondo, M., Gennaro, R.: Provably secure threshold password-authenticated key exchange.

- Shared secret key between two or more entities
- Take advantages of the distributed system distributed authentication
 - Robustness, scalability and greater availability
- Authentication:
 - Password-based
 - Key agreement and key confirmation between the parties
 - Provably secure protocol
 - Efficiency

Proposed protocol

I J_{ν} Servers $(K_1, \ldots, K_k), G$ K_{ν}, G $\overline{K}_1, \dots, \overline{K}_{k-1}$ short-lived keys $K_i = KKDF_{kev}^{c+i}(psw)$, where key = H(salt||psw) $K_n = KKDF_{key}(psw) \oplus \cdots \oplus KKDF^{c+n-2}_{kev}(psw)$ $t_1, \ldots, t_{k-1}, t_v; r_1, \ldots, r_{k-1}, x random$ $w_1 = H(t_1), \ldots, w_v = H(t_v)$ $w = H(w_1 || \dots || w_{k-1} || w_v)$ $m_0 = H(w)$ $m_i = (Mac_{K_i}(r_i \oplus J_i) \oplus w_i) ||r_i|$ $m_v = (Mac_{K_v}(r_v \oplus xG \oplus J_v) \oplus w_v)||r_v||xG$ $M_1{=}I||J_1||...||J_k||m_0||...||m_k$ public channel $I || m_i$ J_i K_i, \overline{K}_i $m_i = p || q$

 $w'_i = p \oplus Mac_{K_i}(q \oplus J_i)$

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Proposed protocol

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Servers

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$$\begin{array}{c} Enc_{\overline{K}_{1}}\left(w_{1}^{\prime}\right)\\ \\ \hline m_{r}=u||s||z\\ w_{r}^{\prime}=u\oplus Mac_{K_{1}}(s\oplus z\oplus J_{r})\\ w^{\prime}=H(w^{\prime})||...||w_{1}^{\prime}..||w_{r}^{\prime}\\ \\ m_{0}=H(w)^{2}=H(w^{\prime})\\ \\ y \ random value\\ \\ ssk=H_{0}(yxG)\\ \\ h=H(ssk||sG||xG||w^{\prime}) \end{array}$$

 $M_2 = h || yG$

public channel

 $ssk' = H_0(yxG)$

 $h \stackrel{?}{=} H(ssk'||yG||xG||w)$

$$M_3 = H(ssk||yG||xG)$$

 $M_3 \stackrel{?}{=} H(ssk||yG||xG)$

public channel

Security goals

- Correctness
- Key secrecy
- Known-key security (Freshness)
- Mutual authentication
- (Perfect) Forward-secrecy

Bellare-Rogaway model - Indirect proof

- Each participant is modelled by an oracle
- Oracles keep transcripts and answer the questions on the tape in one step.
- It is assumed that the attacker is able to create queries. Each query models some type of attack.

Theorem

The proposed protocol is a secure AKC protocol in the random oracle model, assuming MAC is existentially unforgeable under an adaptive chosen-message attack and symmetric encryption scheme is indistinguishable under chosen plaintext attack, moreover ECCDH assumption holds in the elliptic curve group.

Proof Consider an adversary $\mathcal A$ and suppose that

$$Pr[\text{No-Matching}^{\mathcal{A}}(\kappa)]$$

is non-negligible. There are two cases: either the edge or the client oracle is accepted.

Security assumption of the symmetric encryption

- $n_C(\kappa)$ indicates the probability of an event that the attacker is successful
- Suppose that a client oracle is accepted a server oracle is impersonated by the attacker
- Generating an \mathcal{F} polynomial-time algorithm to break the symmetric encryption scheme is *indistinguishable under chosen plaintext attack*

$$\xi_2(\kappa) = \frac{n_C(\kappa)}{T_1(\kappa)T_2(\kappa)\binom{T_2(\kappa)-1}{k-1}T_3(\kappa)} - \lambda(\kappa),$$

• It contradicts the security assumption of the symmetric encryption



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